

# **Semidefinite Relaxation for a Class of Robust QCQPs: A Verifiable Sufficient Condition for Rank-One Solutions**

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# Problem Statement

**Problem:** A robust quadratically constrained quadratic program (QCQP)

$$\begin{aligned} \min_{\mathbf{w}_1, \dots, \mathbf{w}_K \in \mathbb{C}^N} \quad & \sum_{i=1}^K \|\mathbf{w}_i\|_2^2 \\ \text{s.t.} \quad & \max_{\|\mathbf{h}_i - \bar{\mathbf{h}}_i\|_2 \leq \varepsilon_i} \sigma_i^2 + \mathbf{h}_i^H \left( \sum_{j \neq i} \mathbf{w}_j \mathbf{w}_j^H - \frac{1}{\gamma_i} \mathbf{w}_i \mathbf{w}_i^H \right) \mathbf{h}_i \leq 0, \quad i = 1, \dots, K, \end{aligned}$$

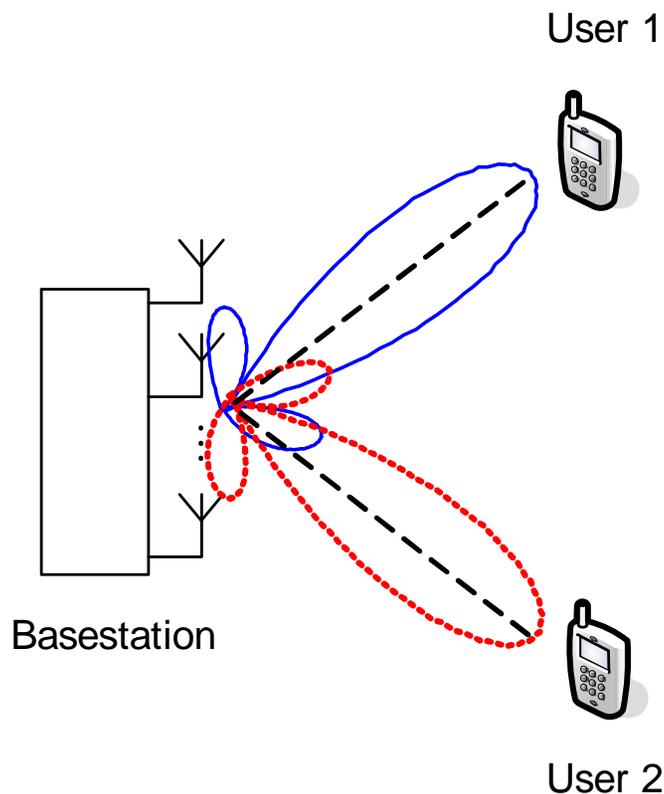
where  $\bar{\mathbf{h}}_i \in \mathbb{C}^N$ ,  $\sigma_i^2, \varepsilon_i, \gamma_i > 0$ ,  $i = 1, \dots, K$ , are given.

- non-convex
- may be approximated by techniques like convex restrictions and relaxations

**Question:** How well does [semidefinite relaxation \(SDR\)](#) perform?

# Motivating Application: Downlink Beamforming in Communications

**Scenario:** a base station (BS) sending  $K$  independent information signals to  $K$  users simultaneously; BS has  $N$  transmit antennas; users have one receive antenna.



**Quality-of-service characterization:** the signal-to-interference-and-noise ratios (SINRs)

$$\text{SINR}_i = \frac{|\mathbf{h}_i^H \mathbf{w}_i|^2}{\sum_{j \neq i} |\mathbf{h}_i^H \mathbf{w}_j|^2 + \sigma_i^2}, \quad i = 1, \dots, K,$$

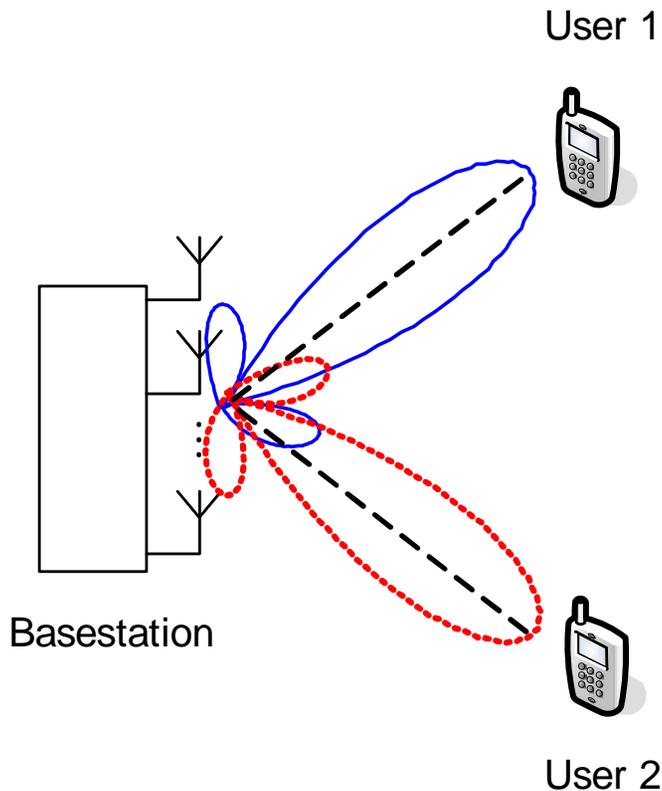
where

$\mathbf{h}_i \in \mathbb{C}^N$  is the channel from the BS to user  $i$ ;

$\mathbf{w}_i \in \mathbb{C}^N$  the beamforming vector of user  $i$ ;

$\sigma_i^2$  the noise power.

# A Downlink Beamforming Formulation



**Problem:** an SINR-constrained design

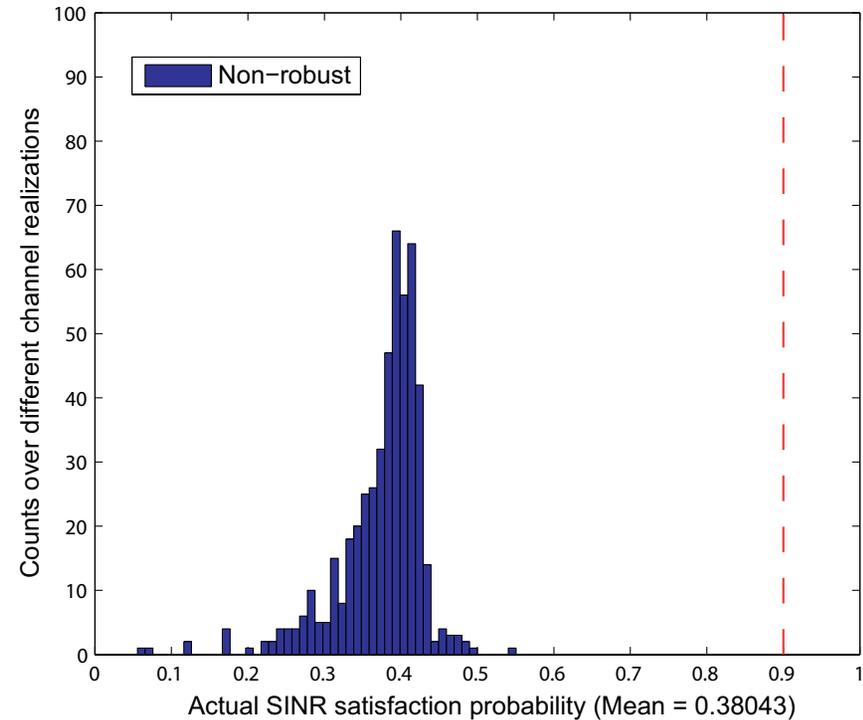
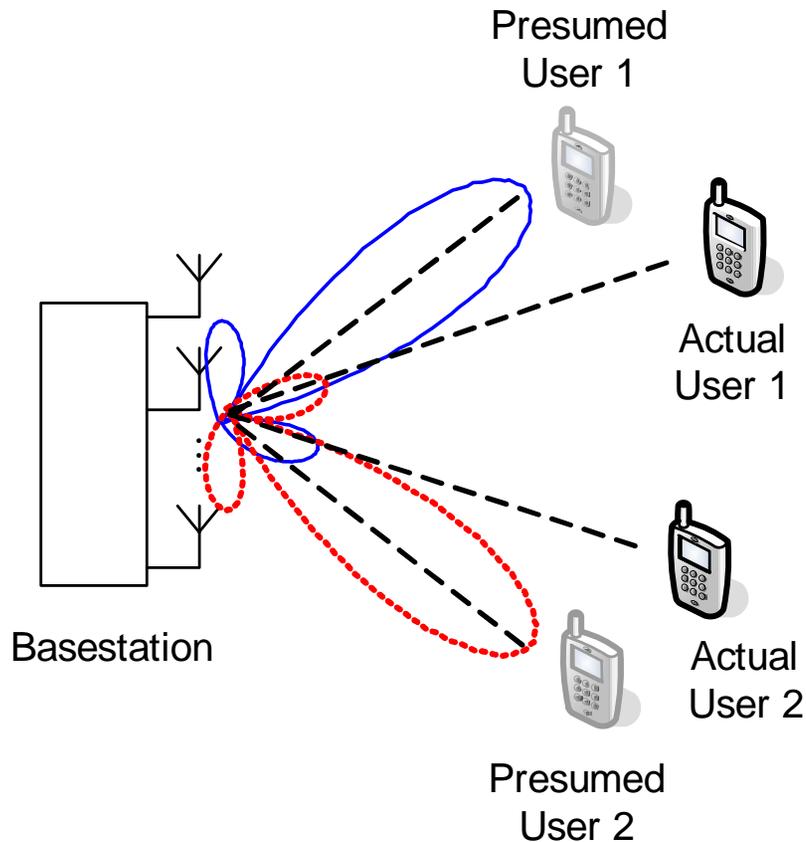
$$\begin{aligned} \min_{\mathbf{w}_1, \dots, \mathbf{w}_K \in \mathbb{C}^N} \quad & \sum_{i=1}^K \|\mathbf{w}_i\|_2^2 \\ \text{s.t.} \quad & \text{SINR}_i = \frac{|\mathbf{h}_i^H \mathbf{w}_i|^2}{\sum_{j \neq i} |\mathbf{h}_i^H \mathbf{w}_j|^2 + \sigma_i^2} \geq \gamma_i, \\ & i = 1, \dots, K, \end{aligned}$$

where  $\gamma_i$  is user- $i$ 's minimum SINR requirement.

# Sensitivity Issues under Imperfect Channel Information

## Issue:

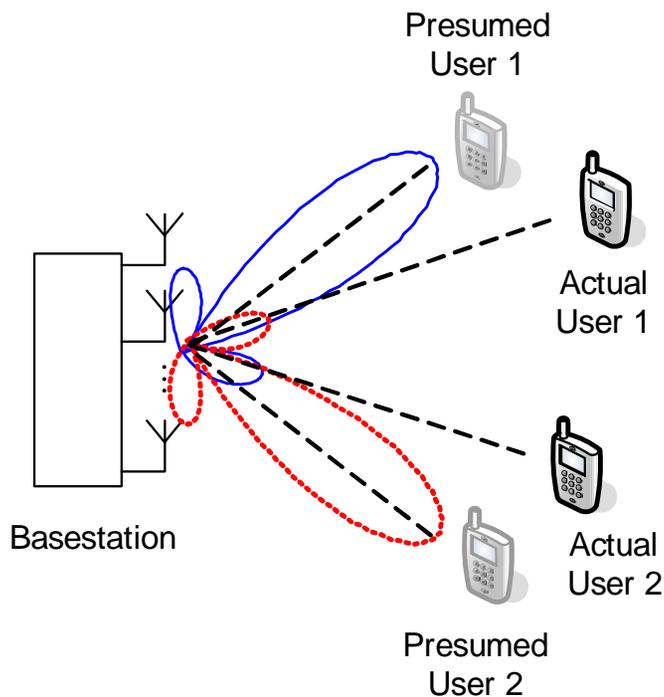
- the SINR-constrained design assumes that the channels  $h_1, \dots, h_K$  are perfectly known at the BS;
- in practice,  $h_1, \dots, h_K$  are often imperfectly known



More than 50% SINR outage!

# Robustifying the Beamforming Design

**Goal:** Guarantee that the SINR requirements are satisfied under any spherically bounded channel uncertainties.



**Problem:**

$$\begin{aligned} \min_{\mathbf{w}_1, \dots, \mathbf{w}_K \in \mathbb{C}^N} \quad & \sum_{i=1}^K \|\mathbf{w}_i\|_2^2 \\ \text{s.t.} \quad & \min_{\|\mathbf{h}_i - \bar{\mathbf{h}}_i\|_2 \leq \varepsilon_i} \text{SINR}_i \geq \gamma_i, \\ & i = 1, \dots, K, \end{aligned}$$

where  $\bar{\mathbf{h}}_i$  is the presumed channel of user  $i$ ;  $\varepsilon_i$  is the uncertainty radius.

## A Review of SDR: The Non-Robust Case

Recall the (non-robust) SINR-constrained design

$$\begin{aligned} \min_{\mathbf{w}_1, \dots, \mathbf{w}_K \in \mathbb{C}^N} \quad & \sum_{i=1}^K \|\mathbf{w}_i\|_2^2 \\ \text{s.t.} \quad & \text{SINR}_i = \frac{|\mathbf{h}_i^H \mathbf{w}_i|^2}{\sum_{j \neq i} |\mathbf{h}_i^H \mathbf{w}_j|^2 + \sigma_i^2} \geq \gamma_i, \quad i = 1, \dots, K. \end{aligned}$$

**SDR:** apply  $\mathbf{W}_i = \mathbf{w}_i \mathbf{w}_i^H \iff \mathbf{W}_i \succeq \mathbf{0}, \text{rank}(\mathbf{W}_i) \leq 1$  to the above problem, and drop the rank constraints to obtain a relaxed problem

$$\begin{aligned} \min_{\mathbf{W}_1, \dots, \mathbf{W}_K \succeq \mathbf{0}} \quad & \sum_{i=1}^K \text{Tr}(\mathbf{W}_i) \\ \text{s.t.} \quad & \sigma_i^2 + \mathbf{h}_i^H \left( \sum_{j \neq i} \mathbf{W}_j - \frac{1}{\gamma_i} \mathbf{W}_i \right) \mathbf{h}_i \leq 0, \quad i = 1, \dots, K. \end{aligned}$$

- convex, a semidefinite program (SDP)
- **Question:** Is SDR tight? Or, does SDR always admit a rank-one solution?

# Rank-One Solution Guarantee via SDP Rank Reduction

Consider an extension of the Shapiro-Barvinok-Pataki (SBP) rank reduction result:

**Fact [Huang-Palomar'09]:** Consider a complex-valued SDP

$$\begin{aligned} \min_{\mathbf{W}_1, \dots, \mathbf{W}_k \succeq \mathbf{0}} \quad & \sum_{i=1}^k \text{Tr}(\mathbf{C}_i \mathbf{W}_i) \\ \text{s.t.} \quad & \sum_{l=1}^k \text{Tr}(\mathbf{A}_{i,l} \mathbf{W}_l) \geq b_i, \quad i = 1, \dots, m. \end{aligned}$$

If  $m \leq k + 2$  and some mild assumptions hold, then there exists a solution  $(\mathbf{W}_1^*, \dots, \mathbf{W}_k^*)$  such that  $\text{rank}(\mathbf{W}_i^*) = 1$  for all  $i$ .

- SDR is tight for the SINR-constrained problem since  $k = m = K$
- **note:** the same conclusion can also be drawn via other proof approaches, such as uplink-downlink duality [Bengtsson-Ottersten'01] and a “folklore” result (to be explained).

## A Review of SDR: The Robust Case

The SDR of the robust SINR-constrained design:

$$\min_{\mathbf{W}_i \succeq \mathbf{0} \forall i} \sum_{i=1}^K \text{Tr}(\mathbf{W}_i) \quad (\text{P.1})$$

$$\text{s.t.} \quad \max_{\|\mathbf{h}_i - \bar{\mathbf{h}}_i\|_2 \leq \varepsilon_i} \sigma_i^2 + \mathbf{h}_i^H \left( \sum_{j \neq i} \mathbf{W}_j - \frac{1}{\gamma_i} \mathbf{W}_i \right) \mathbf{h}_i \leq 0, i = 1, \dots, K. \quad (\text{P.2})$$

- convex, but (P.2) are semi-infinite
- By applying the  $\mathcal{S}$ -lemma to (P.2), Problem (P) can be reformulated as an SDP

$$\begin{aligned} & \min_{\substack{\mathbf{W}_1, \dots, \mathbf{W}_K \succeq \mathbf{0}, \\ t_1, \dots, t_K \geq 0}} \sum_{i=1}^K \text{Tr}(\mathbf{W}_i) \\ & \text{s.t.} \quad \begin{bmatrix} \mathbf{Q}_i + t_i \mathbf{I} & \mathbf{r}_i \\ \mathbf{r}_i^H & s_i - t_i \varepsilon_i^2 \end{bmatrix} \succeq \mathbf{0}, \quad i = 1, \dots, K, \end{aligned}$$

where  $\mathbf{Q}_i = \frac{1}{\gamma_i} \mathbf{W}_i - \sum_{j \neq i} \mathbf{W}_j$ ,  $\mathbf{r}_i = \mathbf{Q}_i \bar{\mathbf{h}}_i$ ,  $s_i = \bar{\mathbf{h}}_i^H \mathbf{Q}_i \bar{\mathbf{h}}_i - \sigma_i^2$ .

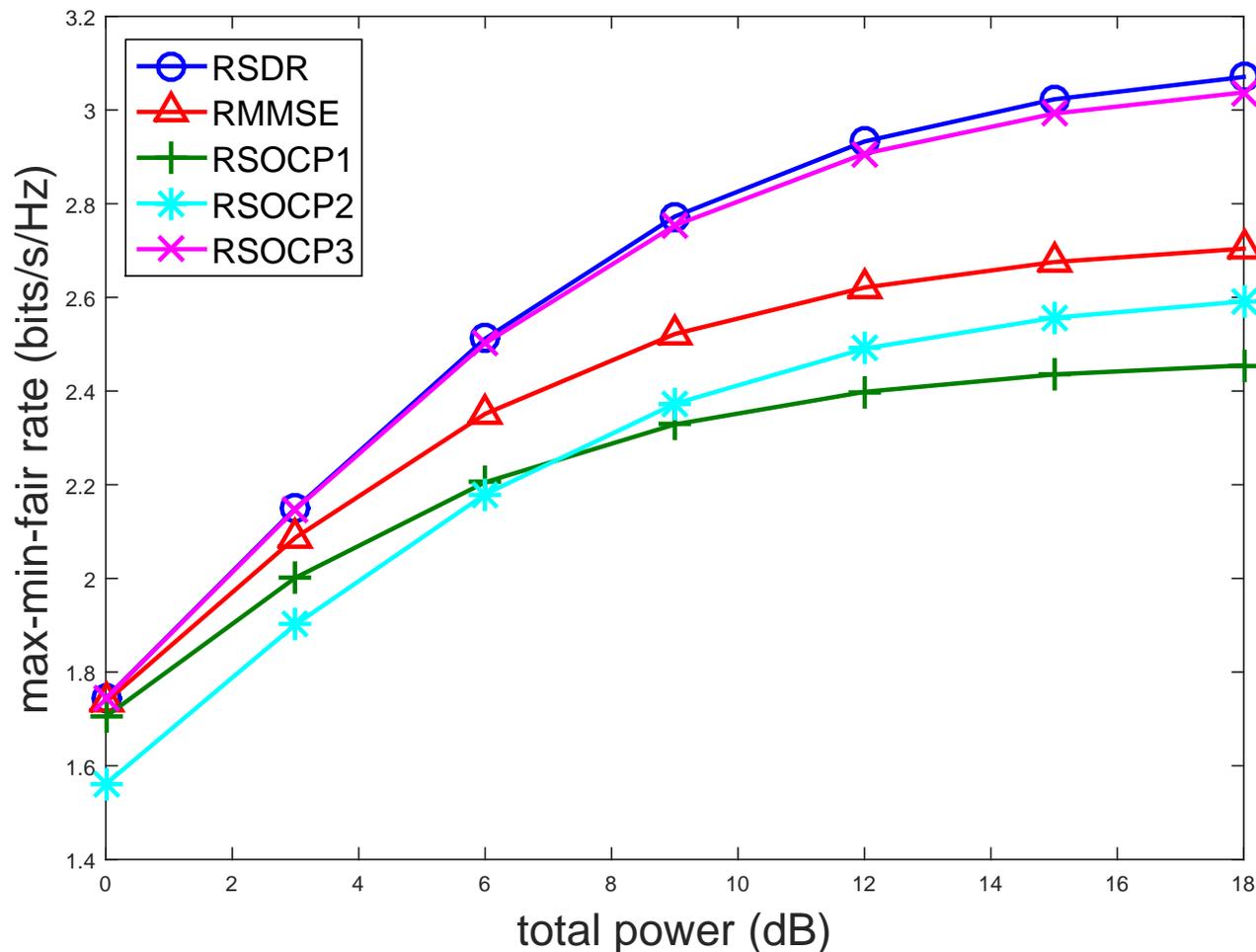
– first proposed in **[Zheng-Wang-Ng'08]**

# A Curious Numerical Finding

**Observation:** The SDR problem was empirically found to admit a rank-one solution in almost all feasible instances!

$r$ (bits/s/Hz)	number of rank-1 instances / number of feasible instances										
	$(N, K) = (4, 3)$		$(N, K) = (8, 3)$		$(N, K) = (8, 7)$		$(N, K) = (12, 7)$		$(N, K) = (12, 11)$		
	$\varepsilon_i^2 = 0.1$	$\varepsilon_i^2 = 0.05$	$\varepsilon_i^2 = 0.1$	$\varepsilon_i^2 = 0.05$	$\varepsilon_i^2 = 0.1$	$\varepsilon_i^2 = 0.05$	$\varepsilon_i^2 = 0.1$	$\varepsilon_i^2 = 0.05$	$\varepsilon_i^2 = 0.1$	$\varepsilon_i^2 = 0.05$	
0.1375	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000
0.2122	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000
0.3233	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000
0.4835	<b>1999/2000</b>	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000
0.7057	<b>1999/2000</b>	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000
1.0000	1973/1973	1995/1995	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000
1.3701	1933/1933	1993/1993	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000
1.8122	1688/1688	1889/1889	2000/2000	2000/2000	<b>1950/1952</b>	1997/1997	2000/2000	2000/2000	2000/2000	2000/2000	2000/2000
2.3165	1535/1535	1833/1833	2000/2000	2000/2000	1084/1084	1814/1814	1999/1999	2000/2000	<b>1483/1485</b>	1976/1976	1976/1976
2.8698	1258/1258	1743/1743	2000/2000	2000/2000	271/ 271	995/ 995	1964/1964	1998/1998	109/ 109	1068/1068	1068/1068
3.4594	839/ 839	1539/1539	1994/1994	2000/2000	51/ 51	549/ 549	1795/1795	1993/1993	6/ 6	160/ 160	160/ 160
4.0746	365/ 365	1187/1187	1961/1961	2000/2000	4/ 4	181/ 181	1262/1262	1936/1936	0/ 0	28/ 28	28/ 28
4.7070	68/ 68	688/ 688	1753/1753	1987/1987	0/ 0	19/ 19	354/ 354	1659/1659	0/ 0	2/ 2	2/ 2
5.3509	1/ 1	211/ 211	955/ 955	1920/1920	0/ 0	0/ 0	12/ 12	885/ 885	0/ 0	0/ 0	0/ 0
6.0022	0/ 0	21/ 21	106/ 106	1485/1485	0/ 0	0/ 0	0/ 0	122/ 122	0/ 0	0/ 0	0/ 0
6.6582	0/ 0	0/ 0	1/ 1	469/ 469	0/ 0	0/ 0	0/ 0	0/ 0	0/ 0	0/ 0	0/ 0

# Comparison with Other Approximation Methods



$N = 4$ ,  $K = 3$ ,  $\sigma_i^2 = 0.1$ ,  $\varepsilon_i^2 = 0.1$ . RSDR= robust SDR; RMMSE= [Vučić-Boche'09], SOCP1= [Shenouda-Davidson'07], SOCP2= [Tajer-Prasad-Wang'11], SOCP3= [Huang-Palomar-Zhang'13]. The benchmarked methods are convex restrictions.

# Our Main Interest

The SDR problem:

$$\begin{aligned} \min_{\mathbf{W}_i \succeq \mathbf{0} \ \forall i} \quad & \sum_{i=1}^K \text{Tr}(\mathbf{W}_i) \\ \text{s.t.} \quad & \max_{\|\mathbf{h}_i - \bar{\mathbf{h}}_i\|_2 \leq \varepsilon_i} \sigma_i^2 + \mathbf{h}_i^H \left( \sum_{j \neq i} \mathbf{W}_j - \frac{1}{\gamma_i} \mathbf{W}_i \right) \mathbf{h}_i \leq 0, \quad i = 1, \dots, K. \end{aligned} \quad (\text{P})$$

**Challenge:** Can we theoretically identify conditions under which Problem (P) is guaranteed to admit a rank-one solution?

# Can We Call Our Old Friend, SBP Rank Reduction?

Recall the SDP form of Problem (P):

$$\begin{aligned} \min_{\substack{\mathbf{W}_1, \dots, \mathbf{W}_K \succeq \mathbf{0}, \\ \mathbf{Z}_1, \dots, \mathbf{Z}_K \succeq \mathbf{0}, \\ t_1, \dots, t_K \geq 0}} \sum_{i=1}^K \text{Tr}(\mathbf{W}_i) \end{aligned} \quad (\text{P-SDP.1})$$

$$\text{s.t. } \mathbf{Z}_i = \begin{bmatrix} \mathbf{Q}_i + t_i \mathbf{I} & \mathbf{r}_i \\ \mathbf{r}_i^H & s_i - t_i \varepsilon_i^2 \end{bmatrix}, \quad i = 1, \dots, K. \quad (\text{P-SDP.2})$$

**Question:** Can we apply SBP rank reduction to Problem (P-SDP), just as in the non-robust case, to obtain a rank-one solution result?

- **Answer:** **No**, at least by our experience.
  - Why? Each matrix equality constraint in (P-SDP.P2) contains many scalar equality constraints.

# An Existing Result by Song, Shi, Sanjabi, Sun and Luo

Denote the optimal value of Problem (P) by

$$v^* = \min_{\mathbf{W}_i, \mathbf{Z}_i \succeq \mathbf{0}, t_i \geq 0, \forall i} \sum_{i=1}^K \text{Tr}(\mathbf{W}_i)$$
$$\text{s.t. } \mathbf{Z}_i = \begin{bmatrix} \mathbf{Q}_i + t_i \mathbf{I} & \mathbf{r}_i \\ \mathbf{r}_i^H & s_i - t_i \varepsilon_i^2 \end{bmatrix}, \quad i = 1, \dots, K.$$

**Result [Song-Shi-Sanjabi-Sun-Luo'12]:** A solution  $(\mathbf{W}_1^*, \dots, \mathbf{W}_K^*)$  to Problem (P) must be of rank one if

$$\varepsilon_i^2 < \frac{\gamma_i \sigma_i^2}{v^*}, \quad \text{for } i = 1, \dots, K.$$

**Implication:** Problem (P) should admit a rank-one solution under sufficiently small error bounds  $\varepsilon_i$ 's.

**Drawback:** **unverifiable;**  $v^*$  also depends on the problem instance  $\{\bar{\mathbf{h}}_i, \sigma_i^2, \varepsilon_i, \gamma_i\}_{i=1}^K$ .

## A Verifiable Result by Us

Let

$$\hat{\mathbf{F}} = [ \bar{\mathbf{h}}_1 / \|\bar{\mathbf{h}}_1\|_2, \dots, \bar{\mathbf{h}}_K / \|\bar{\mathbf{h}}_K\|_2 ]$$

be the presumed multiuser channel direction matrix.

**Result [Ma-Pan-So-Chang'16]:** Under a few mild assumptions, a solution  $(\mathbf{W}_i^*)_{i=1}^K$  to Problem (P) must be of rank one if

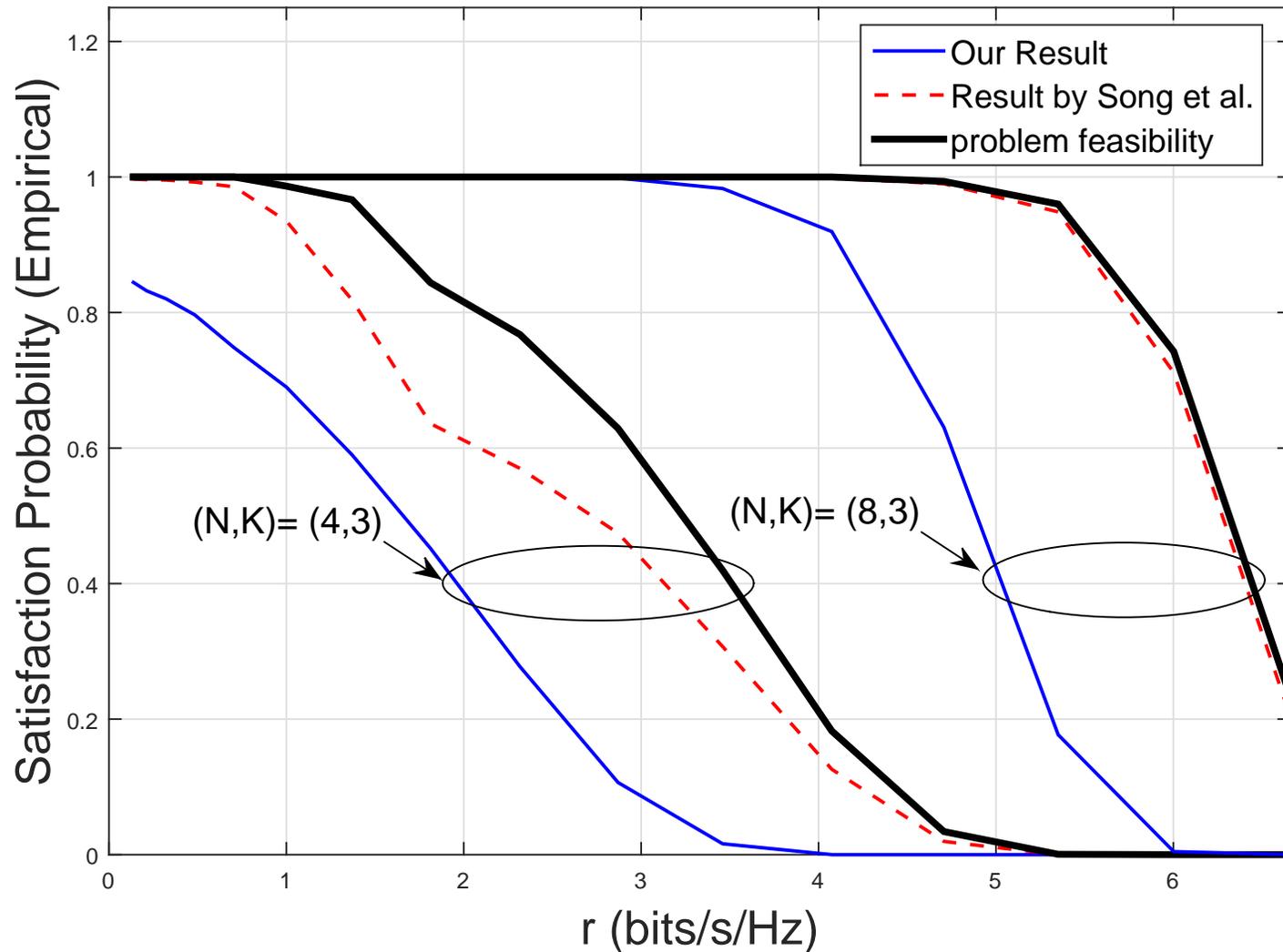
$$\frac{\|\bar{\mathbf{h}}_k\|_2^2}{\varepsilon_k^2} \sigma_{\min}(\hat{\mathbf{F}})^2 > 1 + K + \left( K - \frac{1}{K} \right) \gamma_k, \quad k = 1, \dots, K,$$

where  $\sigma_{\min}(\hat{\mathbf{F}})$  is the smallest singular value of  $\hat{\mathbf{F}}$ .

**Implication:** The SDR problem will admit a rank-one solution if

- the channel-to-uncertainty ratios  $\|\bar{\mathbf{h}}_k\|_2^2 / \varepsilon_k^2$  are sufficiently large;
- the channel direction matrix  $\hat{\mathbf{F}}$  is sufficiently well-conditioned.

# Tightness of Our Verifiable Condition



There is a gap between our verifiable condition and numerical result. Nevertheless, the performance trends of the two are consistent.

# Proof Sketch of Our Result: Setting the Stage

Let us write the robust SDR problem as

$$\begin{aligned} \min_{\mathbf{W} \in \mathcal{S}} \quad & \sum_{i=1}^K \text{Tr}(\mathbf{W}_i) \\ \text{s.t.} \quad & \max_{\|\mathbf{h}_i - \bar{\mathbf{h}}_i\|_2 \leq \varepsilon_i} \varphi_i(\mathbf{W}, \mathbf{h}_i) \leq 0, \quad i = 1, \dots, K. \end{aligned}$$

where  $\mathbf{W} = (\mathbf{W}_1, \dots, \mathbf{W}_K)$ ,  $\mathcal{S} = \{\mathbf{W} \mid \mathbf{W}_i \succeq \mathbf{0} \ \forall i\}$ ,

$$\varphi_i(\mathbf{W}, \mathbf{h}_i) = \sigma_i^2 + \mathbf{h}_i^H \left( \sum_{j \neq i} \mathbf{W}_j - \frac{1}{\gamma_i} \mathbf{W}_i \right) \mathbf{h}_i.$$

- $\varphi_i$  is affine in  $\mathbf{W}$  and **indefinite** in  $\mathbf{h}_i$

# Proof Sketch of Our Result: An Equivalent Representation of the Robust Constraints

Let's do SDR with the robust constraint functions:

$$\begin{aligned}
 & \max_{\|\mathbf{h}_i - \bar{\mathbf{h}}_i\|_2 \leq \varepsilon_i} \varphi_i(\mathcal{W}, \mathbf{h}_i) \\
 &= \max_{\|\mathbf{h}_i - \bar{\mathbf{h}}_i\|_2 \leq \varepsilon_i} \sigma_i^2 + \text{Tr} \left( \mathbf{h}_i \mathbf{h}_i^H \left( \sum_{j \neq i} \mathbf{W}_j - \frac{1}{\gamma_i} \mathbf{W}_i \right) \right) \\
 &\leq \max_{\mathbf{H}_i \in \mathcal{V}_i} \sigma_i^2 + \text{Tr} \left( \mathbf{H}_i \left( \sum_{j \neq i} \mathbf{W}_j - \frac{1}{\gamma_i} \mathbf{W}_i \right) \right) \triangleq \phi_i(\mathcal{W}, \mathbf{H}_i)
 \end{aligned}$$

where  $\mathcal{V}_i = \{\mathbf{H}_i \mid \exists \mathbf{h}_i \text{ s.t. } \mathbf{H}_i \succeq \mathbf{h}_i \mathbf{h}_i^H, \|\bar{\mathbf{h}}_i\|_2^2 - 2\text{Re}(\bar{\mathbf{h}}_i^H \mathbf{h}_i) + \text{Tr}(\mathbf{H}_i) \leq \varepsilon_i^2\}$ .

SDR is tight in this case (SBP rank reduction). Thus,

$$\max_{\|\mathbf{h}_i - \bar{\mathbf{h}}_i\|_2 \leq \varepsilon_i} \varphi_i(\mathcal{W}, \mathbf{h}_i) = \max_{\mathbf{H}_i \in \mathcal{V}_i} \phi_i(\mathcal{W}, \mathbf{H}_i).$$

- $\phi_i$  is affine in  $\mathcal{W}$  and affine in  $\mathbf{H}_i$ .

# Proof Sketch of Our Result: A New Duality Result

**Theorem:** Under a few mild assumptions, we have

$$\begin{aligned} \min_{\mathcal{W} \in \mathcal{S}} \sum_i \text{Tr}(\mathbf{W}_i) &= \max_{\mathcal{H}_i \in \mathcal{V}_i \forall i} \min_{\mathcal{W} \in \mathcal{S}} \sum_i \text{Tr}(\mathbf{W}_i) \\ \text{s.t. } \max_{\mathcal{H}_i \in \mathcal{V}_i} \phi_i(\mathcal{W}, \mathbf{H}_i) &\leq 0, \forall i & \text{s.t. } \phi_i(\mathcal{W}, \mathbf{H}_i) &\leq 0, \forall i \end{aligned}$$

Also,

$\mathcal{W}^*$  is a solution to the LHS problem  $\implies$  there exists  $\mathcal{H}^*$  such that  $(\mathcal{H}^*, \mathcal{W}^*)$  is a maximin solution to the RHS problem

- Proof Idea: Sion's maximin theorem and some simple arguments; the affine property of  $\phi_i$  is crucial.

## Proof Sketch of The New Duality Theorem

$$\begin{aligned}
 \text{LHS Problem} &= \min_{\mathbf{W} \in \bar{\mathcal{S}}} \sup_{\boldsymbol{\lambda} \geq \mathbf{0}} \sum_i \text{Tr}(\mathbf{W}_i) + \sum_i \lambda_i \sup_{\mathbf{H}_i \in \mathcal{V}_i} \phi_i(\mathcal{W}, \mathbf{H}_i) \\
 &= \sup_{\boldsymbol{\lambda} \geq \mathbf{0}} \min_{\mathbf{W} \in \bar{\mathcal{S}}} \sum_i \text{Tr}(\mathbf{W}_i) + \sum_i \lambda_i \sup_{\mathbf{H}_i \in \mathcal{V}_i} \phi_i(\mathcal{W}, \mathbf{H}_i) & \text{(a)} \\
 &= \sup_{\boldsymbol{\lambda} \geq \mathbf{0}} \min_{\mathbf{W} \in \bar{\mathcal{S}}} \sup_{\mathbf{H}_i \in \mathcal{V}_i \forall i} \sum_i \text{Tr}(\mathbf{W}_i) + \sum_i \lambda_i \phi_i(\mathcal{W}, \mathbf{H}_i) \\
 &= \sup_{\boldsymbol{\lambda} \geq \mathbf{0}} \sup_{\mathbf{H}_i \in \mathcal{V}_i \forall i} \min_{\mathbf{W} \in \bar{\mathcal{S}}} \sum_i \text{Tr}(\mathbf{W}_i) + \sum_i \lambda_i \phi_i(\mathcal{W}, \mathbf{H}_i) & \text{(b)} \\
 &= \sup_{\mathbf{H}_i \in \mathcal{V}_i \forall i} \sup_{\boldsymbol{\lambda} \geq \mathbf{0}} \min_{\mathbf{W} \in \bar{\mathcal{S}}} \sum_i \text{Tr}(\mathbf{W}_i) + \sum_i \lambda_i \phi_i(\mathcal{W}, \mathbf{H}_i) \\
 &= \sup_{\mathbf{H}_i \in \mathcal{V}_i \forall i} \min_{\mathbf{W} \in \bar{\mathcal{S}}} \sup_{\boldsymbol{\lambda} \geq \mathbf{0}} \sum_i \text{Tr}(\mathbf{W}_i) + \sum_i \lambda_i \phi_i(\mathcal{W}, \mathbf{H}_i) & \text{(c)} \\
 &= \text{RHS Problem}
 \end{aligned}$$

where (a), (b) and (c) are all due to Sion's maximin theorem.

- all about flipping min and sup!
- the affine property of  $\phi_i$  is essential in (b).

# Proof Sketch of Our Result: Further Discussion

**Theorem:** Under a few mild assumptions, we have

$$\begin{aligned} \min_{\mathcal{W} \in \mathcal{S}} \sum_i \text{Tr}(\mathbf{W}_i) &= \max_{\mathcal{H}_i \in \mathcal{V}_i} \min_{\mathcal{W} \in \mathcal{S}} \sum_i \text{Tr}(\mathbf{W}_i) \\ \text{s.t. } \max_{\mathcal{H}_i \in \mathcal{V}_i} \phi_i(\mathcal{W}, \mathbf{H}_i) &\leq 0, \quad \forall i & \text{s.t. } \phi_i(\mathcal{W}, \mathbf{H}_i) &\leq 0, \quad \forall i \end{aligned}$$

Also,

$$\begin{array}{l} \mathcal{W}^* \text{ is a solution} \\ \text{to the LHS problem} \end{array} \implies \begin{array}{l} \text{there exists } \mathcal{H}^* \text{ such that} \\ (\mathcal{H}^*, \mathcal{W}^*) \text{ is a maximin solution} \\ \text{to the RHS problem} \end{array}$$

Discussion:

- every inner problem on the RHS has a rank-one solution (SBP rank reduction)
- does that imply that the LHS problem has a rank-one solution?

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**Discussion:**

- every inner problem on the RHS has a rank-one solution (SBP rank reduction)
- does that imply that the LHS problem has a rank-one solution?
  - **No**, the theorem didn't say

$$\begin{aligned} (\mathcal{H}^*, \mathcal{W}^*) \text{ is a maximin solution} & \implies \mathcal{W}^* \text{ is a solution} \\ \text{to the RHS problem} & \text{to the LHS problem} \end{aligned}$$

# Proof Sketch of Our Result: Further Discussion

**Theorem:** Under a few mild assumptions, we have

$$\begin{aligned} \min_{\mathcal{W} \in \mathcal{S}} \sum_i \text{Tr}(\mathbf{W}_i) &= \max_{\mathcal{H}_i \in \mathcal{V}_i \forall i} \min_{\mathcal{W} \in \mathcal{S}} \sum_i \text{Tr}(\mathbf{W}_i) \\ \text{s.t. } \max_{\mathcal{H}_i \in \mathcal{V}_i} \phi_i(\mathcal{W}, \mathbf{H}_i) \leq 0, \forall i &\quad \text{s.t. } \phi_i(\mathcal{W}, \mathbf{H}_i) \leq 0, \forall i \end{aligned}$$

Also,

$\mathcal{W}^*$  is a solution to the LHS problem  $\implies$  there exists  $\mathcal{H}^*$  such that  $(\mathcal{H}^*, \mathcal{W}^*)$  is a maximin solution to the RHS problem

**Discussion:**

- however, if every inner problem on the RHS **must** admit a rank-one solution, then the solution  $\mathcal{W}^*$  to the LHS problem must be of rank one.
  - why don't we check when such instances happen?

# Proof Sketch of Our Result: A Different Rank-One Result

Consider the non-robust SINR-constrained design

$$\begin{aligned} \min_{\mathbf{W} \in \mathcal{S}} \quad & \sum_i \text{Tr}(\mathbf{W}_i) \\ \text{s.t.} \quad & \phi_i(\mathbf{W}, \mathbf{H}_i) = \sigma_i^2 + \text{Tr} \left( \mathbf{H}_i \left( \sum_{j \neq i} \mathbf{W}_j - \frac{1}{\gamma_i} \mathbf{W}_i \right) \right) \leq 0, \quad i = 1, \dots, K \end{aligned} \tag{P2}$$

where  $\mathbf{H}_i \in \mathcal{V}_i$  for all  $i$ .

**Aim:** Identify conditions under which a solution to (P2) **must have** rank one.

- SBP rank reduction doesn't work; it's only good at saying "there exists"

## Proof Sketch of Our Result: A Different Rank-One Result

Consider the non-robust SINR-constrained design

$$\begin{aligned} & \min_{\mathbf{W} \in \mathcal{S}} \sum_i \text{Tr}(\mathbf{W}_i) \\ & \text{s.t. } \phi_i(\mathbf{W}, \mathbf{H}_i) = \sigma_i^2 + \text{Tr} \left( \mathbf{H}_i \left( \sum_{j \neq i} \mathbf{W}_j - \frac{1}{\gamma_i} \mathbf{W}_i \right) \right) \leq 0, \quad i = 1, \dots, K \end{aligned} \quad (\text{P2})$$

where  $\mathbf{H}_i \in \mathcal{V}_i$  for all  $i$ .

**Fact** (folklore): If all  $\mathbf{H}_i$ 's take a rank-one form  $\mathbf{H}_i = \mathbf{h}_i \mathbf{h}_i^H$ , then a solution to (P2) must have rank one.

- Proof Idea: exploit the specific structures of the dual of (P2). Particularly, the dual variables of (P2) w.r.t.  $\mathbf{W}_i$ 's take the form

$$\mathbf{Z}_i = \mathbf{I} + \sum_{j \neq i} \mu_j \mathbf{H}_j - \frac{\mu_i}{\gamma_i} \mathbf{H}_i \succeq \mathbf{0}, \quad \boldsymbol{\mu} \geq \mathbf{0} \implies \text{rank}(\mathbf{Z}_i) \geq N - 1$$

The complementary slackness  $\mathbf{Z}_i \mathbf{W}_i = \mathbf{0}$  enforces  $\text{rank}(\mathbf{W}_i) \leq 1$ .

# Proof Sketch of Our Result: A Different Rank-One Result

Consider the non-robust SINR-constrained design

$$\begin{aligned} \min_{\mathbf{W} \in \mathcal{S}} \quad & \sum_i \text{Tr}(\mathbf{W}_i) \\ \text{s.t.} \quad & \phi_i(\mathbf{W}, \mathbf{H}_i) = \sigma_i^2 + \text{Tr} \left( \mathbf{H}_i \left( \sum_{j \neq i} \mathbf{W}_j - \frac{1}{\gamma_i} \mathbf{W}_i \right) \right) \leq 0, \quad i = 1, \dots, K \end{aligned} \tag{P2}$$

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**Fact** (folklore): If all  $\mathbf{H}_i$ 's take a rank-one form  $\mathbf{H}_i = \mathbf{h}_i \mathbf{h}_i^H$ , then a solution to (P2) must have rank one.

## Our Finishing Touch:

- every  $\mathbf{H}_i \in \mathcal{V}_i$  can be written as  $\mathbf{H}_i = \mathbf{h}_i \mathbf{h}_i^H + \mathbf{\Xi}_i$  for some  $\mathbf{h}_i$ ,  $\mathbf{\Xi} \succeq \mathbf{0}$ ;
- study a variation of the folklore fact for  $\mathbf{H}_i = \mathbf{h}_i \mathbf{h}_i^H + \mathbf{\Xi}_i$  (with  $\mathbf{\Xi}_i$  being small);
- identify conditions under which (P2) must have rank-one solutions for all  $\mathbf{H}_i \in \mathcal{V}_i$

## Conclusion and Discussion

- We considered a specific robust QCQP and showed a verifiable sufficient condition under which SDR is tight.
- Future challenge: Can we establish a strong rank-one solution result? Simulation results indicate the SDR solution is almost always of rank one.

**Thank you.** Preprint available on <https://arxiv.org/abs/1602.01569> or



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